

Analysis and Comparison of Embedded Network Stacks

Design and Evaluation of the GNRC Network Stack

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- 1. Introduction
- 2. RIOT
- 3. GNRC
- 4. Evaluation of GNRC
- 5. Conclusion



1. Introduction

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The Internet of Things

- Freie Universität
- "Internet of Things" = broad, generic term
 - Home automation
 - Industry 4.0
 - Cars
 - Health surveillance
 - Wildlife surveillance
 - Wireless sensor networks
 - ...



The IoT – Constraints & Requirements



- Large address space: > 10 Internet connected devices per person
- Low energy requirements
 - Low processing power: a few MHz
 - Small memory: \leq 10 KiB RAM, \leq 100 KiB flash
 - Lossy transmission medium: IEEE 802.15.4, Bluetooth Low-Energy, NFC

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- Constraints govern need for:
 - specific OSs: TinyOS, Contiki, FreeRTOS, RIOT
 - specific communication protocols: ZigBee, Z-Wave, IETF's IPv6-based IoT suite



Approaching a solution

Problem 1 Large address space



 $2^{32}\approx 4.3\cdot 10^9 {\rm possible} ~{\rm addresses} \ll 7.4\cdot 10^{10} {\rm devices}$

 \Rightarrow **IPv6** (2¹²⁸ $\approx 3.4 \cdot 10^{38}$)



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 \Rightarrow IPv6 (2¹²⁸ $\approx 3.4 \cdot 10^{38}$)?

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Problem 4 No TCP, no HTTP, no WWW? \Rightarrow Non-TCP alternative **CoAP**



Application	HTTP	DNS		CoAP	DNS					
Transport	ТСР	UDP		U	ЭР					
Network	IPv4 /	/ IPv6								
Data link	Link	Layer		link Laver						
Physical										

Traditional TCP/IP stack

IoT stack by IETF



Existing solutions

 $\begin{array}{l} \mbox{existing stack (RIOT)} \\ (+) \mbox{ IoT support} \end{array}$



existing stack (RIOT)

- (+) IoT support
- (-) Very rigid in selection of protocols
- (-) Single-packet buffering
- (-) no clear structure / unmaintainable



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 \Rightarrow Multi-interface support required (only BLIP and IwIP provides that)

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Functional Requirements:

- Focus on IoT protocols
- Multiple interface support
- Ability to handle >1 packet at a time

Non-functional Requirements:

- Open Standards and Tools
- Comprehensive configurability
- Modularity
- Low Memory Footprint (< 10 KiB RAM, < 30 KiB code-size)
- Low-Power Design



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Scheduler:

- Tick-less scheduling policy (O(1)):
 - Highest priority thread runs until finished or blocked
 - ISR can preempt any thread at all time
 - If all threads are blocked or finished:
 - Special IDLE thread is run
 - Goes into low-power mode



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IPC:

Synchronous (default) and asynchronous (optional, by IPC queue initialization)

RIOT's Networking architecture

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- devised to integrate any network stack into RIOT



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Common network device API:

netdev_t		netdev_driver_t
+event_callback : callback_t +context : void*	* 1	+send(data_buf) : int +recv(data buf) : int
+send(data_buf) : int +recv(data_buf) : int +get(opt_type, opt_buf) : int	-	<pre>+get(opt_type, opt_buf) : int +set(opt_type, opt_buf) : int +isr() : void</pre>
<pre>+set(opt_type, opt_buf) : int +isr() : void</pre>	_	

isr() method allows for getting out of ISR context









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- collection of unified connectivity APIs to the transport layer
- What's the problem with POSIX sockets?
 - too generic for most use-cases
 - numerical file descriptors (internal storage of state required)
 - in general: too complex for usage, too complex for porting
- protocol-specific APIs:
 - conn_ip (raw IP)
 - conn_udp (UDP)
 - conn_tcp (TCP)
 - ...
- both IPv4 and IPv6 supported



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The components of GNRC







netapi

- Inter-modular API utilizing IPC
- Two asynchronous message types (don't expect reply) for data transfer:
 - GNRC_NETAPI_MSG_TYPE_SND: pass "down" the stack (send)
 - GNRC_NETAPI_MSG_TYPE_RCV: pass "up" the stack (receive)
- Two synchronous message types (expect reply) for option handling:
 - GNRC_NETAPI_MSG_TYPE_GET: get option value
 - GNRC_NETAPI_MSG_TYPE_SET: set option value
- specification deliberately vague

 \Rightarrow implementations can make own preconditions on data



Network interfaces in GNRC (1)





Network interfaces in GNRC (1)





Network interfaces in GNRC (2)



- netapi-capable thread as any other protocol implementation
- implement MAC protocol
- communication to driver via netdev
 - \leftarrow timing requirements for e.g. TDMA-based MAC protocols





How to know where to send netapi messages?



- How to know where to send netapi messages?
- Both protocol implementation and users can register to be interested in type + certain context (e.g. port in UDP)
 - gnrc_netreg_register(GNRC_NETTYPE_IPV6, ALL, &me)
 - gnrc_netreg_register(GNRC_NETTYPE_UDP, PORT_DNS, &me)
- \Rightarrow Find handler for packets in registry

pktbuf



- Data packet stored in pktbuf
- Representation: list of variable-length "packet snips"
- Protocols can *mark* sections of data to create new snip
- keeping track of referencing threads: reference counter users
 - if users == 0: packet removed from packet buffer
 - if users > 1 and write access requested: packet duplicated (copy-on-write)
- to keep duplication minimal: only up to current snip
 - \Rightarrow Reverse order of snips (not data) on reception





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Feature-based comparison

 Comparison of GNRC with emb6 (OS-independent fork of uIP) and IwIP

			6L c	W P	WPAN			ICMPv6										
	ace.	Frag.												RPL				
Stack	multi-ifa	r eseq.	mult.	HC1	IPHC	NHC	IPv6	error	echo	NDP	SLAAC	6Lo-ND	MLD	st.	non-st.	тсь	UDP	C oA P
GNRC IwIP	~	×	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	× ×	~	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	~	•	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	~	×	×	×	×	×	×	~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	*
emb6	×	~	×	•	~	•	•	~	•	~	×	×	×	~	×	•	~	~

Comparison of network stack features (✓ = supported, X = not supported, ● = partially supported, ◆ = support through external library)

IwIP additionally has IPv4 (+ ARP), PPP and DNS support



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Stack-traversal time





Memory usage



- Taken application for stack traversal time tests in reception as reference
- compiled on 32-bit platform (ARM Cortex-M3)
- network stacks were configured to handle 1280 byte IPv6 packets





- overall close second behind lwIP
- considering GNRC's age (~1 yr vs. ~15 yr of lwIP and uIP) ⇒ very good
- GNRC easier to work with
 - configuration of both emb6 and lwIP fiddly
 - documentation: mixed reactions from community

Discussion of GNRC



Advantages

- Well defined interface enforces clear communication between modules
- Use-cases are easy to describe in terms of API usage
- IPC-based API allows parallel data handling per design
- Very loose coupling between modules
- packet buffer's size easy to adapt to given use-case

Disadvantages

- IPC-based API is hard to debug
- memory hungry due to required memory stack allocation
- theory vs. praxis: cross-layer requirements everywhere



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Contributions

- Co-design of GNRC, netdev, and conn
- Implementation work:
 - over 500 PRs contributed on GitHub:
 - 6LoWPAN and IPv6 (incl. NDP) layer for GNRC
 - pktbuf
 - several netdev-based drivers
 - port of lwIP and emb6 to RIOT
 - ...
- RIOT maintenance:
 - over 500 PRs (co-)reviewed on GitHub
 - consultance to community regarding all things GNRC
- Research:
 - co-authorship and presentation of paper to workshop @ ACM MobiSys'15
 - co-authorship of proposed paper to USENIX OSDI'16



- Performance-wise GNRC only (close) second after more mature lwIP
- BUT: GNRC developed with real-time in mind, lwIP not
- Both GNRC and emb6 can be stripped down via configuration to be smaller
- GNRC remains best candidate for embedded RTOS RIOT





- Optimization efforts both size- and performance-wise
- Mitigation efforts of GNRC's disadvantages
- Expansion of GNRC's feature set.
- Further experimentation with other testing parameters
 - power consumption
 - performance under stress
 - ...
- Further experimentation with more stacks
 - BLIP (TinyOS)
 - vanilla uIP and RIME (Contiki)
 - OpenWSN
 - CCN-lite

• ...